

# SIFA: Exploiting Ineffective Fault Inductions on Symmetric Cryptography

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Robert Primas<sup>1</sup>

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<sup>2</sup>Infineon Technologies AG, Germany  
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We present fault attacks that are ...

- Hard to prevent
  - Defy detection, any degree of redundancy
  - Defy infection
  - (Defy masking)
- Versatile
  - Many possible fault locations/effects
  - Applicable to many symmetric schemes
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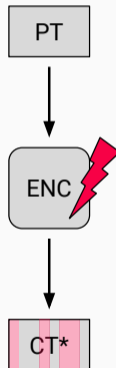
# Fault Attacks

- Get device access:
    - Set plaintexts
    - Observe ciphertexts
  - Cause (partially) erroneous computation
  - Observe faulty and correct ciphertext
  - Determine correct sub key guesses by verifying output pairs
- ⇒ Differential Fault Attack (DFA)



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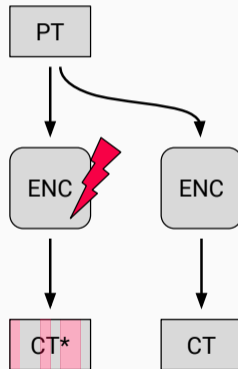
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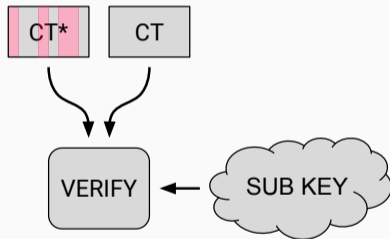
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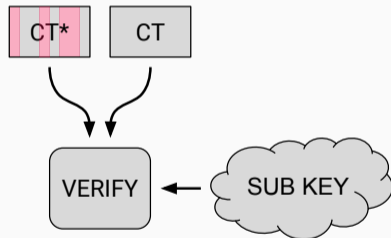




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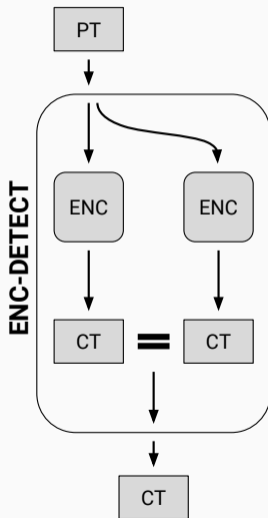
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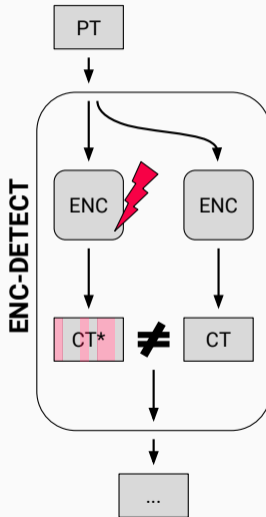
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- Use redundancy to detect faults
  - Fault detected → No ciphertext
  - 2 identical faults necessary for attack
- More redundancy, Enc-Dec, masking, etc...



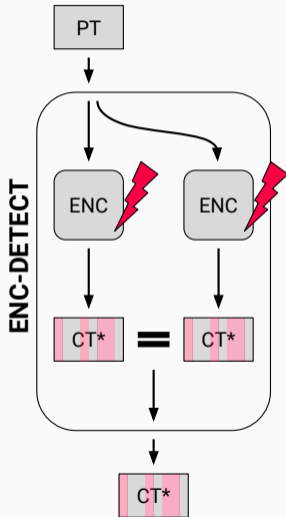
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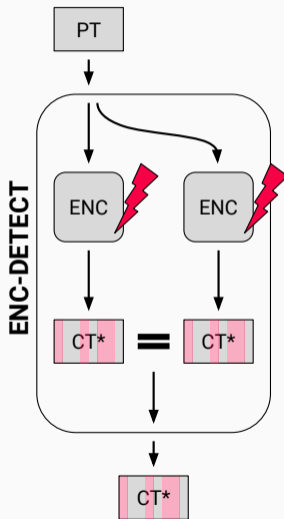
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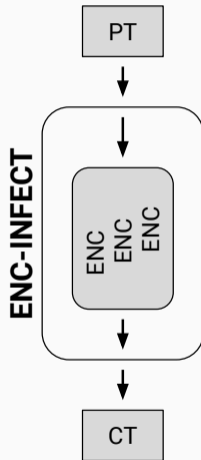
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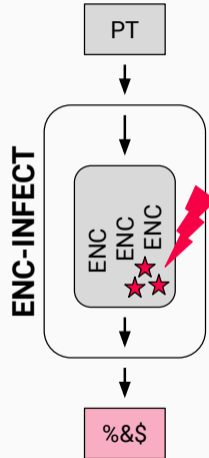
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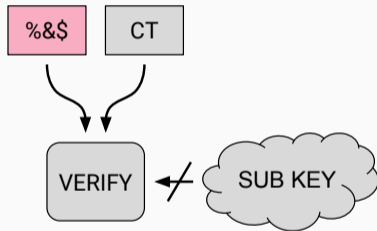
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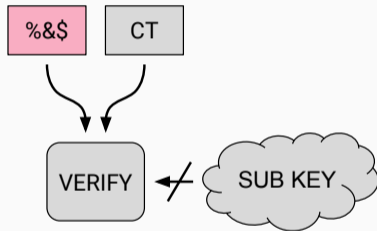
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Combines ...

- Ineffective Fault Attacks (IFA) by Clavier et al. [Cla07]
  - + Exploits only correct ciphertexts (similar to safe error attacks)
  - Requires precise faults with known effect
- Statistical Fault Analysis (SFA) by Fuhr et al. [FJLT13]
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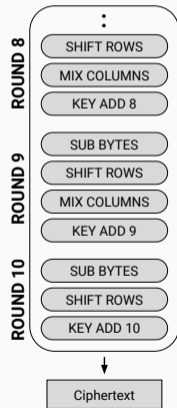
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# SIFA on AES - Fault Injection Phase

## Example for AES...

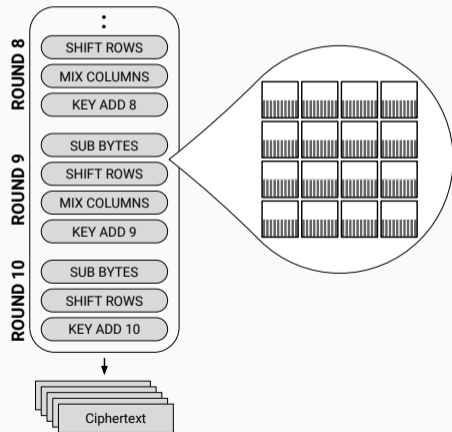
- Over multiple encryptions, state bytes are uniformly distributed
- Fault somewhere between MC in round 8-9
- Goal is some non-uniform distribution
  - Stuck-at fault, random fault, skips, flips...
  - Fault Granularity: 1 bit → a few bytes
- **Works even for ineffective faults**
  - i.e. a fault was injected but the computation is still correct
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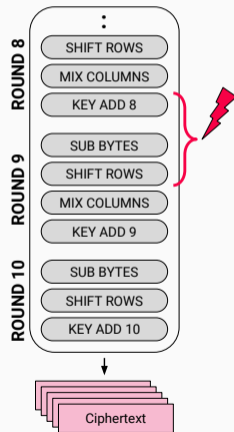
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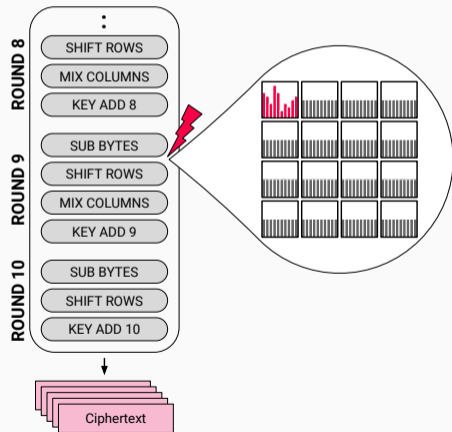
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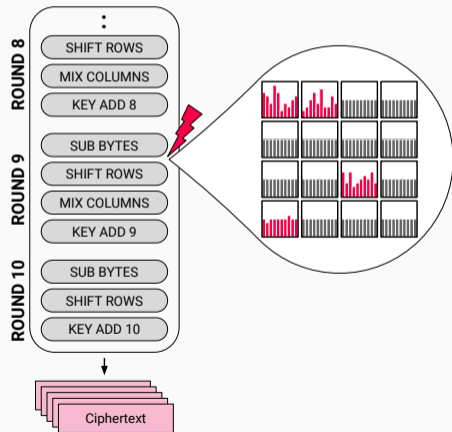




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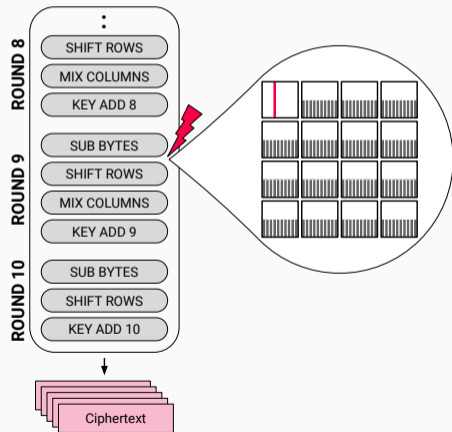
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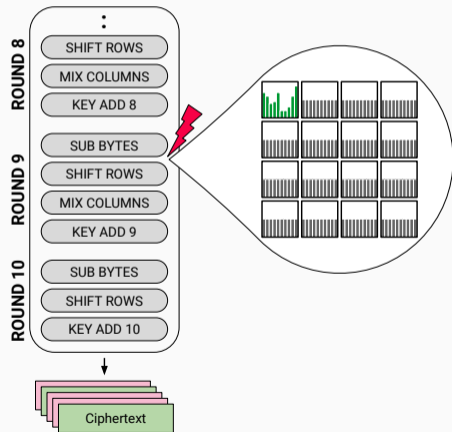
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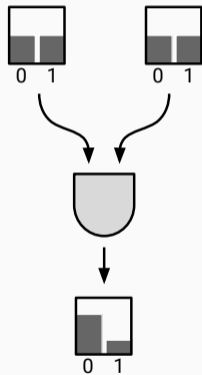
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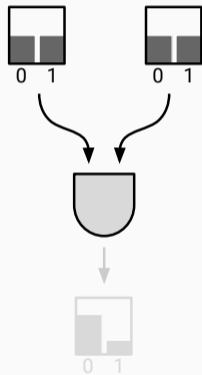
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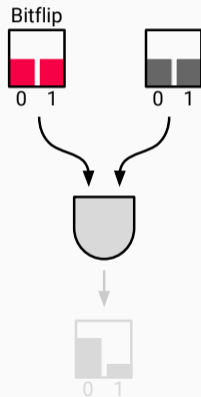
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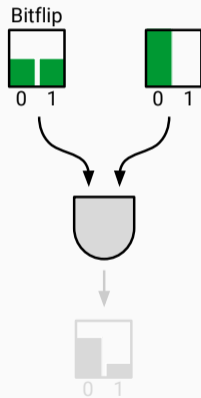
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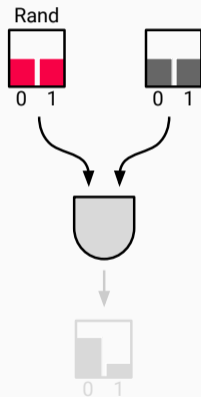
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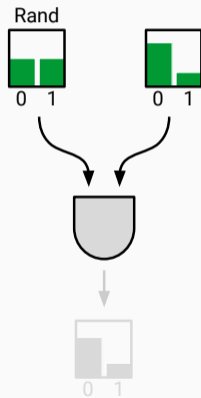


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## SIFA on AES - Key Recovery Phase

- Collect set of correct ciphertexts  $\mathcal{C}_1 \dots \mathcal{C}_n$  from faulted encryptions
- Guess 32-bit sub key  $\mathcal{K}_{10}$  and calculate state  $\mathcal{S}_i$  in round 9 ( $\mathcal{K}_9$  is not needed):

$$\mathcal{S}_i = \text{MC}^{-1} \circ \text{SB}^{-1} \circ \text{SR}^{-1}(\mathcal{C}_i \oplus \mathcal{K}_{10})$$

- Measure uniformity of  $\mathcal{S}_1 \dots \mathcal{S}_n$  using e.g. the Squared Euclidean Imbalance (SEI)
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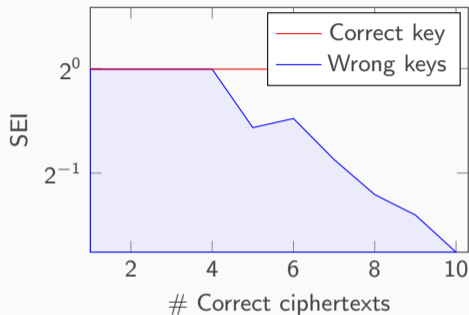
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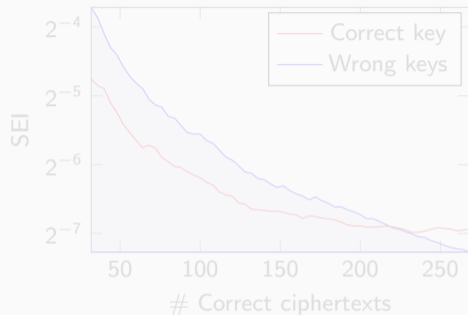
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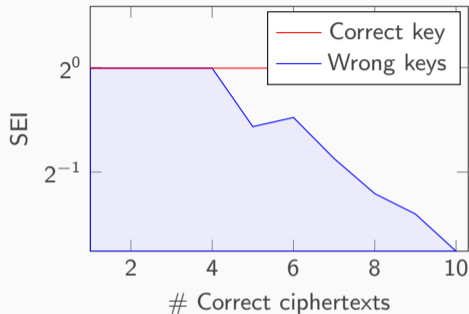
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- SW-AES from AVR-crypto-lib
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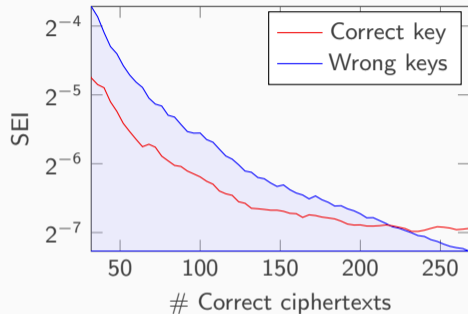
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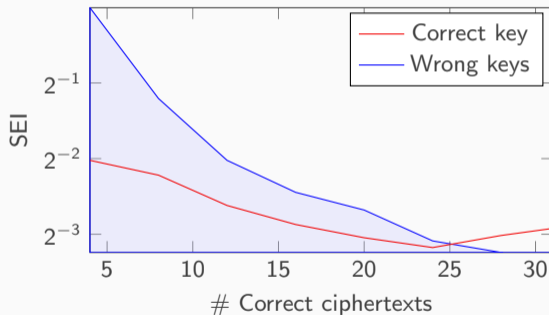
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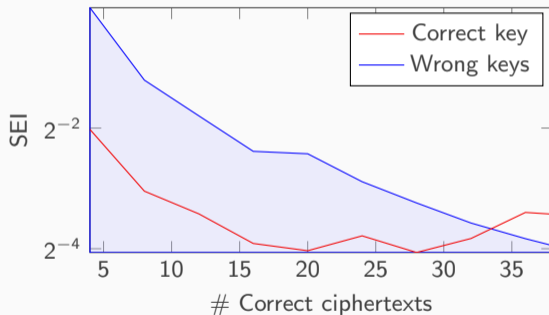
## Results - Infection by Tupsamudre et al. [TBM14]

- Clock glitch: ATXmega128D4
- SW-AES with infection
- 22 real + **11** dummy rounds
- $\approx$  **25** correct ciphertexts
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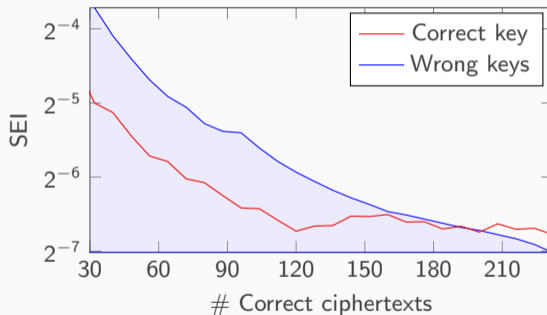
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- $\approx$  **9 000** faulted encryptions



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- Clock glitch: ATXmega128D4
- SW-AES with infection
- 22 real + **66** dummy rounds
- $\approx$  **180** ciphertexts needed
- $\approx$  **46 000** faulted encryptions



## SIFA ...

- defies popular fault countermeasures: detection/infection
- requires hundreds/thousands faulted computations
- requires only one fault per computation
- does not require precise fault locations
- works with any type of fault, even if effect is unknown (→ blackbox attacks)

⇒ works for AE schemes (SAC 2018) [DMMP18]

→ including stream-cipher, sponge-based schemes

→ e.g. all CAESAR finalists

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→ same performance, no real overhead

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**Thank you for your attention!**



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**Secret external encodings do not prevent transient fault analysis.**  
In Pascal Paillier and Ingrid Verbauwhede, editors, *Cryptographic Hardware and Embedded Systems – CHES 2007*, volume 4727 of *LNCS*, pages 181–194.  
Springer, 2007.
-  Christoph Dobraunig, Maria Eichlseder, Hannes Gross, Stefan Mangard, Florian Mendel, and Robert Primas.  
**Statistical ineffective fault attacks on masked AES with fault countermeasures.**  
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**Fault attacks on nonce-based authenticated encryption: Application to keyak and ketje.**

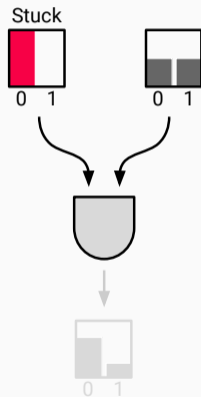
To appear at: Selected Areas of Cryptography, 2018.

 Thomas Fuhr, Éliane Jaulmes, Victor Lomné, and Adrian Thillard.  
**Fault attacks on AES with faulty ciphertexts only.**

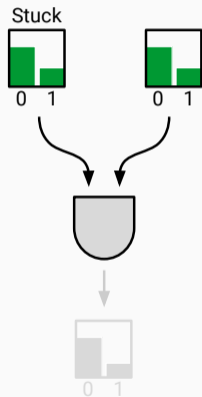
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## SIFA Intuition (cont.)



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